

Linear Programming Primer

Wednesday, 24 August 2022 12:03 AM

A problem of the form: choose x_1, x_2, \dots, x_n
to minimize $c_1 x_1 + \dots + c_n x_n$
subject to: $a_{11} x_1 + a_{12} x_2 + \dots + a_{1n} x_n \geq b_1$
 $a_{21} x_1 + a_{22} x_2 + \dots + a_{2n} x_n \geq b_2$
 \vdots
 $a_{m1} x_1 + \dots + a_{mn} x_n \geq b_m$

where c_i 's, a_{ij} 's, & b_j 's are given constants
(i.e., minimize a linear function of the variables x_1, \dots, x_n
subject to m linear constraints)

is called a **linear program**,

Note that 'min' can be replaced by 'max',

$$\begin{aligned} \geq & \text{ by } \leq \\ = & \text{ by } \leq, \geq \end{aligned}$$

Eg.:

$$\begin{array}{ll} \max & -5x_1 + 2x_2 \\ & x_1 + x_2 \geq 4 \\ & 2x_1 - 5x_2 \geq 0 \\ & -x_1 - x_2 \geq -7 \end{array} \quad \equiv \quad \begin{array}{ll} \min & 5x_1 - 2x_2 \\ & -x_1 - x_2 \leq -4 \\ & 2x_1 - 5x_2 \geq 0 \\ & x_1 + x_2 \leq 7 \end{array}$$

A **feasible solution** to a linear program is an assignment of values to the variables that satisfies all the constraints

An **optimal solution** is feasible, & minimizes the objective among all feasible solutions.

For the LP above, $(4, 0)$ is a feasible solution, w/ objective value -20 .

The optimal value is $(\frac{20}{7}, \frac{8}{7})$, w/ objective value -12 .

An LP can be written in **matrix form**:

$$(P) \quad \min \quad c^T x \quad \text{where } c, x \in \mathbb{R}^n \\ \text{s.t. } A x \geq b \quad A \in \mathbb{R}^{m \times n}, b \in \mathbb{R}^m$$

in the example above, $c = \begin{bmatrix} 5 \\ -2 \end{bmatrix}$, $A = \begin{bmatrix} 1 & 1 \\ 2 & -5 \\ -1 & -1 \end{bmatrix}$, $b = \begin{bmatrix} 4 \\ 0 \\ -7 \end{bmatrix}$

Every LP has a **dual** linear program:

$$(D) \quad \max \quad b^T y \\ \text{s.t. } A^T y = c \\ y \geq 0$$

As can be expected, the dual of (D) is (P).

Strong Duality Theorem: If (P) & (D) are feasible, then the optimal primal value equals the optimal dual value. That is, if x^* & y^* are optimal values for the primal & dual respectively, then
$$c^T x^* = b^T y^*$$

Weak Duality Theorem: If x^* & y^* are feasible values for the primal & dual respectively, then
$$c^T x^* \geq b^T y^*$$

$$\begin{array}{ccc} \hline & | & | \\ & c^T x^* & b^T y^* \\ \hline \text{primal: minimize} & & \text{dual: maximize} \end{array}$$

Finally, using interior-point methods, we can obtain an optimal solution to any LP in polynomial time.